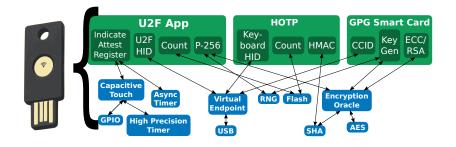
Safely and Efficiently Multiprogramming a 64kB Computer

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Emerging class of embedded applications are software platforms, rather than single purpose devices.

Embedded Software

- ▶ No isolation between components
- Deeply coupled components
- Static memory allocation to avoid unrecoverable runtime memory exhaustion
- ► Fixed concurrency at compile-time

Embedded Hardware

- ► Low-power budget—micro-amps average current consumption
- ▶ 64kB of RAM
- Memory Protection Unit—a limited hardware protection mechanism

Challenges

- How to isolate components despite minimnal hardware resouces?
- ► How to replace individual components without restarting the whole system?
- How to avoid fixed concurrency with limited memory

Common Solutions

► Give up on isolation—write completely bug-free code

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- Whole system updates only

Common Solutions

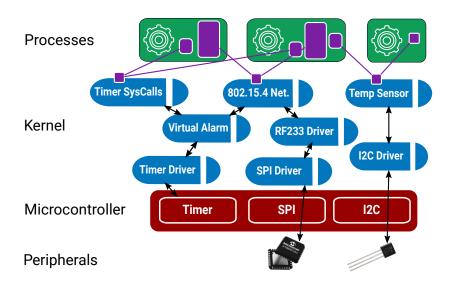
- ► Give up on isolation—write completely bug-free code
- Whole system updates only
- Use *nix et al—forget about low power

Tock

Tock is a new operating system for low-power platforms that takes advantage of the limited hardware-protection mechanisms available on recent microcontrollers and the type-safety features of the Rust programming language to provide a multiprogramming environment:

- Isolation of software faults
- Efficient memory protection and management for dynamic application workloads
- Update/restart/remove individual (user-space) components independently
- ► Retains dependability requirements of long-running devices.

Tock Architecutre



Capsules

- Capsules are components in the kernel
- Minimal runtime overhead:
 - Isolated "at compile-time" using the Rust language type/module system
 - Cooperatively scheduled
 - Can eliminate most isolation at compile-time

Capsules can...

- Violate real-time guarantees
- ▶ Panic (sort of... lets talk...)

But they cannot...

- Read arbitrary memory (secret encryption keys)
- Communicate with peripherals it's not allowed to



Capsules

Stronger memory isolation than hardware protection?

```
struct DMAChannel {
    ...
    enabled: bool,
    buffer: &'static [u8],
}
```

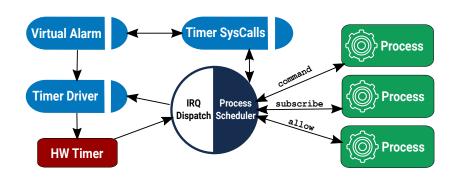
Typing hardware register can constrain allowed values with very fine granularity.

Processes

Can be unreliable since the system can respawn or kill processes without affecting other functionality.

- Hardware isolated concurrent executions of programs
- ▶ Written in any language (currently C, C++, Lua and Rust-ish)
- ► Total control over their memory, including dynamic heap allocation.
- Similar to processes in other systems.
 - Separate stacks allows preemptive execution
 - Memory isolated by the hardware
- Interact with kernel over a small but flexible system-call interface:
 - command, subscribe, allow
 - ▶ yield, memop





What happens when the kernel requires dynamic resources to respond to a request from a process?

- We want to allow arbitrary apps so we don't know concurrency requirements:
 - How many timers will an application need?
 - Will it use SPI, UART, USB, Bluetooth, etc? One socket? 1000 sockets?
- ▶ If the kernel allocates memory for requests dynamically, it may run out of resources.

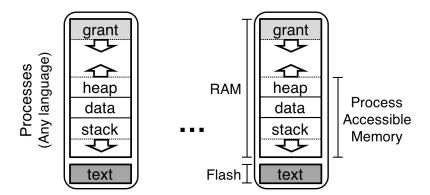
| Threads | Kernel RAM | Syscall RAM | Max Used |
|---------|---------------|----------------|-------------|
| 1 | 3506 | 712 | 158 |
| 2 | 4216 | 1422 | 316 |
| 3 | 4928 | 2134 | 474 |

TOSThreads has low memory efficiency. Static allocation costs 710-712 bytes per thread, of which at most 158 bytes (22%) can be in use at any time. These numbers do not include the thread stacks, each of which can be less than 100 bytes.

Grants

Tock allows a process to "grant" to the kernel portions of its own memory, which the kernel can use to maintain state for process requests.

- Separate sections of kernel heap located in each process's memory space.
- Grant allocations for one process do not affect kernel's ability to allocate for another.
- ► Type-safe interface guarantees all grants for a process can be freed immediately if the process dies.
- ▶ Basic idea: kernel API ensures there are no long-lived pointers directly to grant-allocated memory.



Grant Requirements

- Process cannot access grant allocated memory
 - ▶ We use an additional, dynamically determined MPU rule
- Ensure grant-allocated values unavailable to capsules once process dies through limited API:
 - ► Capsules pass a closure to the enter method
 - ▶ Memory in a grant region only accessible from within closure
 - ▶ Pointers to grant memory cannot escape the closure
 - Implications on kernel design: should avoid cross process data-structures

```
impl<T: Default> Grant {
  fn create() -> Grant<T>
  fn enter<F,R>(&self, proc_id: ProcId, func: F)
    -> Result<R, Error> where
     F: for<'b> FnOnce(&'b mut Owned<T>) -> R, R: Copy
  fn each<F>(&self, func: F) where
     F: for<'b> Fn(&'b mut Owned<T>)
}
```

Grants Compared to the Alternative

Recall: TOSThreads requires 700 bytes statically allocated in the kernel for each additional thread. At most 22% can be used at any given time.

- ► Grants require *no additional per-thread memory* in the kernel
- Only useful memory is dynamically allocated in grants
- Zero wasted memory since it can re-use memory for non-concurrent operations.

Conclusion

- Resource constraints continue to be a challenge for embedded system designers.
 - Low-power, small form-factors and lower cost
- ► These limitations *should not* preclude software abstractions and protections common in general-purpose computers.
- ► Tock provides both dynamic operation and dependability in resource-constrained settings.
 - Best of all: flexible multiprogramming, isolation, system dependability
- Grants split the kernel heap across processes, allowing dynamic demands for kernel resources despite limited system memory

Buy a Hail! https://tockos.org/hardware/hail